

Debts from last week

① "Needle in the haystack" argument:

Fix a tester  $V$ . Pick  $\vec{x}_0 \in \mathbb{F}^m$  uniformly at random.

Let  $p \in \mathbb{F}_{\leq d}[x_1, \dots, x_m]$ ,  $f_{\vec{x}_0}: \mathbb{F}^m \rightarrow \mathbb{F}$

$$f_{\vec{x}_0}(\vec{x}) = \begin{cases} p(\vec{x}) & \vec{x} \neq \vec{x}_0 \\ p(\vec{x}) + 1 & \vec{x} = \vec{x}_0 \end{cases}$$

# of queries performed by  $V$

•  $f_{\vec{x}_0} \notin \mathbb{F}_{\leq d}[x_1, \dots, x_m]$

•  $P_{\vec{x}_0, \pi} [V^{f_{\vec{x}_0, \pi}} \text{ accepts on random } r] \geq P_{\vec{x}_0, r} [V^{p, \pi} \text{ accepts on random } r] - \frac{q}{|\mathbb{F}^m|} \geq 1 - \frac{q}{|\mathbb{F}^m|}$

the proof associated with  $p$

$\Rightarrow$  There exists  $\vec{x}_0 \in \mathbb{F}^m$ , such that

$$\exists \pi \quad P_r [V^{f_{\vec{x}_0, \pi}} \text{ accepts on random } r] \geq 1 - \frac{q}{|\mathbb{F}^m|} > \frac{1}{2}$$

making  $q < |\mathbb{F}^m|/2$  queries to func. and proof

Hence, there is no tester  $V$  sat. the following;

Comp  $p \in \mathbb{F}_{\leq d}[x_1, \dots, x_m] \Rightarrow \exists \pi \quad P_r [V^{p, \pi} \text{ accepts on random } r] = 1$

Sound  $f \notin \mathbb{F}_{\leq d}[x_1, \dots, x_m] \Rightarrow \forall \pi \quad P_r [V^{f, \pi} \text{ accepts on random } r] \leq \frac{1}{2}$

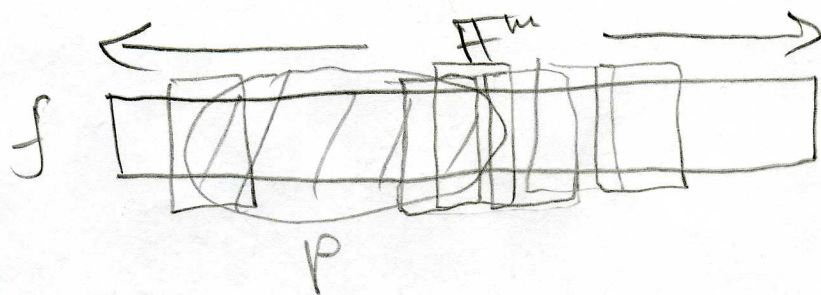
② LDT Theorem  $\Rightarrow$  low degree tester

LDT Theorem For any  $f: \mathbb{F}^m \rightarrow \mathbb{F}$ ,

$$\left| \text{agr}_{\leq d}(f) - \mathbb{E}_{S \in \mathcal{S}_d^m} [\text{agr}_{\leq d}(f|_S)] \right| \leq m^{O(d)} \cdot \left( \frac{d}{|\mathbb{F}|} \right)^{\Omega(d)}$$

$\max_{P \in \mathcal{P}_{\leq d}(\mathbb{F}^m)} \mathbb{P}_{\vec{x}} [f(\vec{x}) = P(\vec{x})]$

all planes in  $\mathbb{F}^m$



Reminder Parameters:

$$d = m \cdot (h-1) \\ = \text{polylog } n$$

$$h^m = n$$

$$h = \log n$$

$$m = \frac{\log n}{\log \log n}$$

$|\mathbb{F}|$  chosen so that  $m^{O(d)} \cdot \left( \frac{d}{|\mathbb{F}|} \right)^{\Omega(d)}$  in the above theorem is small.

$$|\mathbb{F}| = \text{polylog } n.$$

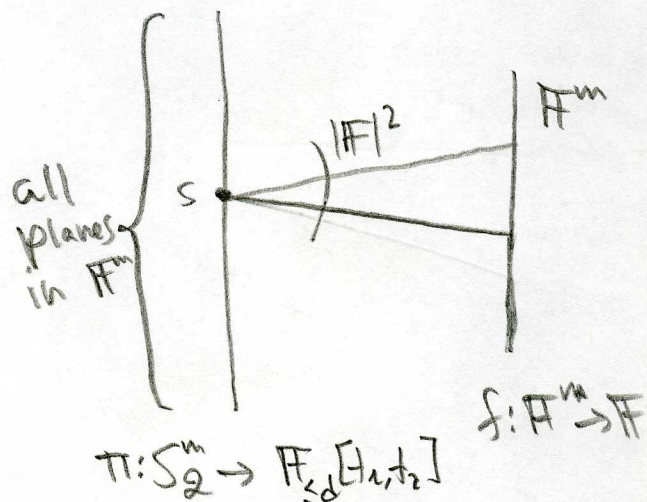
# Plane vs. Point tester

"uniformly at random"

1. Pick u.a.r  $s \in S_2$ ,  $\vec{x} \in S$ .

2. Check  $\pi(s)(t_1, t_2) \stackrel{?}{=} f(\vec{x})$   
 where  $\vec{x} = s(t_1, t_2)$

↑  
 canonical representation of  $s$  as a func.  
 $s: \mathbb{F}^2 \rightarrow \mathbb{F}^m$



## Claim Assuming LDT Theorem,

Comp  $P \in \mathbb{F}_{\leq d}[x_1, \dots, x_m] \Rightarrow \exists \pi$   $P$  (Plane-vs-Point accepts) = 1  
 given access to  $f, \pi$

Sound  $\text{agr}(f, \mathbb{F}_{\leq d}[x_1, \dots, x_m]) \leq 0.8 \Rightarrow \forall \pi$

$P$  (Plane-vs-Point accepts)  $\leq 0.9$   
 given access to  $f, \pi$

Note LDT Thm gives an almost-tight result for all the spectrum.

Pf Comp follows for  $\pi(s) \doteq f|_s$ .

Sound: assuming  $\text{agr}(f, \mathbb{F}_{\leq d}[x_1, \dots, x_m]) \leq 0.8$ . Fix some  $\pi$ .

$$P_{s, \vec{x} \in S} (\text{Plane-vs-Point}^{f, \pi}) \leq \mathbb{E}_{s \in S_2^m} \text{agr}_{\leq d}(f|_s) \leq \text{agr}_{\leq d}(f) + m^{O(d)} \cdot \left(\frac{d}{|\mathbb{F}|}\right)^{\Omega(d)}$$

choosing closest to  $f|_s$  is best possible  $\pi \leq 0.9$   
 ↑  
 for app.  $|\mathbb{F}|$

□

Today Will show:

$$\text{agr}_{\leq d}(f) \geq \mathbb{E}_{S \in S_2^m} [\text{agr}_{\leq d}(f|_S)] - m^{\alpha(d)} \cdot \left(\frac{d}{|F|}\right)^{\Omega(d)}$$

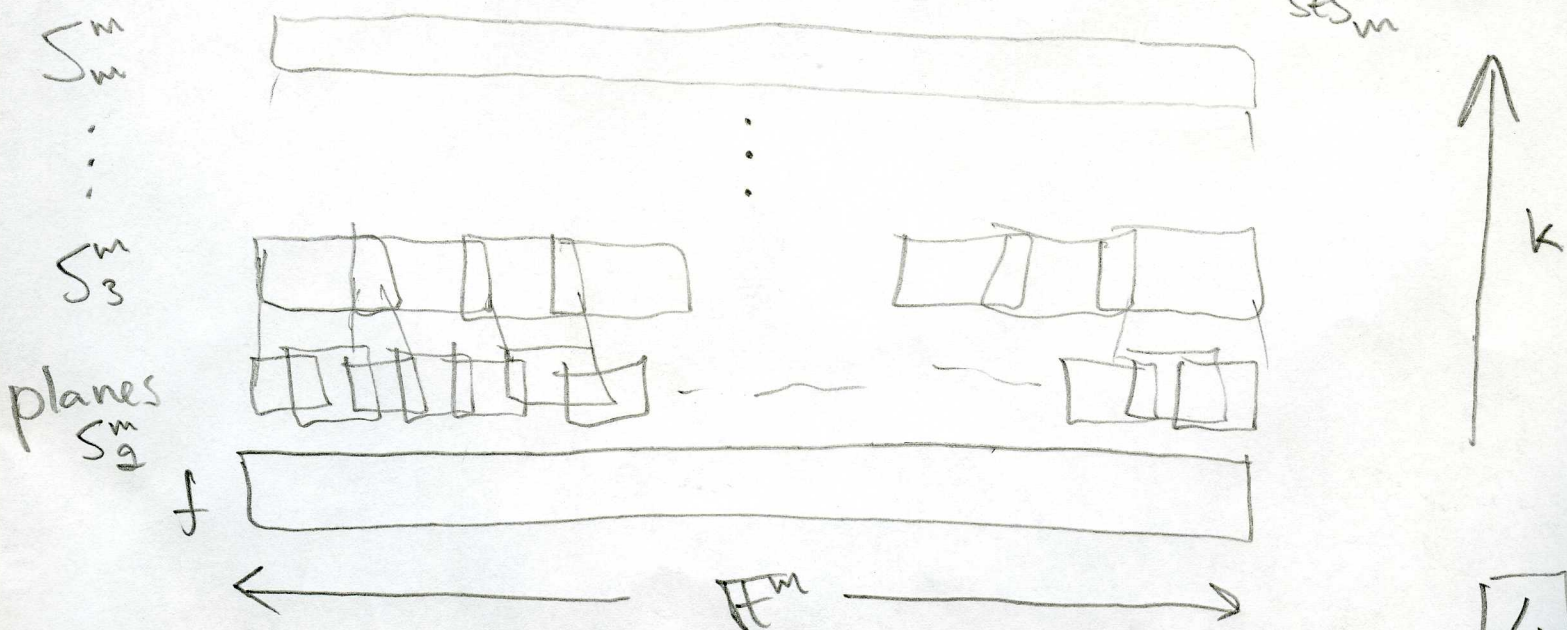
Together with Ex, this will conclude the proof of the LDT Theorem.

(and hence the proof that  $\text{NP} \subseteq \text{PCP}[\mathcal{O}(\log n), \alpha(d)]$  for  $\alpha(d) > 0$ )

Proof strategy Will show by induction that for every  $2 \leq k \leq m$ ,

$$\mathbb{E}_{S \in S_k^m} [\text{agr}_{\leq d}(f|_S)] \geq \mathbb{E}_{S \in S_2^m} [\text{agr}_{\leq d}(f|_S)] - m^{\alpha(d)} \cdot \left(\frac{d}{|F|}\right)^{\Omega(d)}$$

In particular, for  $k=m$ , Note  $\text{agr}_{\leq d}(f) = \mathbb{E}_{S \in S_m^m} [\text{agr}_{\leq d}(f|_S)]$



Important Observation "Symmetry"

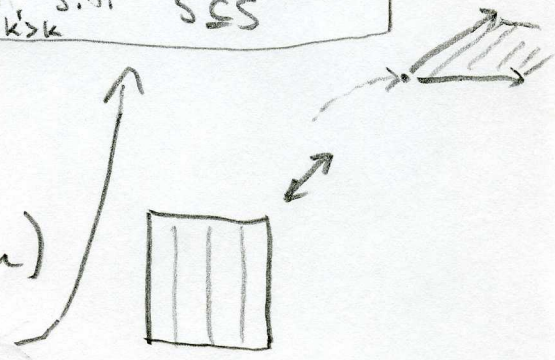
Note for every  $S \in S_k^m$ , there's the same # of  $S' \in S_{k'}^m$  s.t.  $S' \subseteq S$   $k' > k$

For every  $S \in S_{k+1}^m$ ,

there's a bijection  $S \leftrightarrow \mathbb{F}^{k+1}$

(given by the canonical representation)

$$S' \in S_k^m, S' \subseteq S \leftrightarrow S' \in S_k^{k+1}$$



Hence, suffices to prove for  $f: \mathbb{F}^{k+1} \rightarrow \mathbb{F}$ ,

$$(*) \text{ agr}_{\leq d}(f) \geq \mathbb{E}_{S \in S_{k+1}^m} [\text{agr}_{\leq d}(f|_S)] - \epsilon \quad \left[ \text{where } \epsilon = \left( \frac{d}{|\mathbb{F}|} \right)^{2(k+1)} \right]$$

Since

$$(*) \rightarrow \mathbb{E}_{S \in S_k^m} [\text{agr}_{\leq d}(f|_S)] \geq \mathbb{E}_{S \in S_k^m} \left[ \mathbb{E}_{S' \in S_{k+1}^m} [\text{agr}_{\leq d}(f|_{S'})] - \epsilon \right]$$

$$(*) \rightarrow \geq \mathbb{E}_{S \in S_k^m} \left[ \mathbb{E}_{S' \in S_{k+1}^m} \left[ \mathbb{E}_{S'' \in S_{k+1}^m} [\text{agr}_{\leq d}(f|_{S''})] - \epsilon \right] - \epsilon \right]$$

...

$$(*) \rightarrow \geq \mathbb{E}_{S \in S_2^m} [\text{agr}_{\leq d}(f|_S)] - \epsilon m$$

# Bootstrapping argument

We will only prove:

$$\text{agr}_{\leq d}^{(2)}(f) \geq \mathbb{E}_{S \in S_k^{k+1}} [\text{agr}_{\leq d}^{(2)}(f|_S)] - \varepsilon \quad \varepsilon = O\left(\sqrt{\frac{d}{M}}\right)$$

In ex. #3, show this implies

$$\text{agr}_{\leq d}(f) \geq \mathbb{E}_{S \in S_k^{k+1}} [\text{agr}_{\leq d}(f|_S)] - \varepsilon' \quad \varepsilon' \leq \left(\frac{d}{M}\right)^{\Omega(1)}$$

## Hyperplanes Consistency

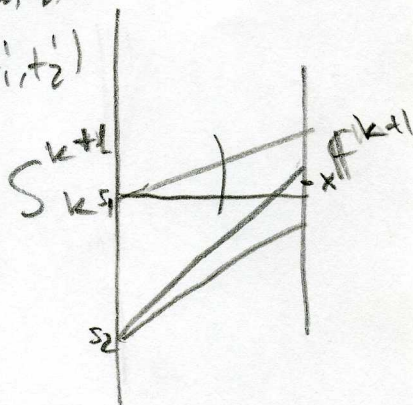
$\pi: S_k^{k+1} \rightarrow \mathbb{F}_{\leq d}(t_1, t_2)$

$$\vec{x} = S_2(t_1, t_2)$$

$$\vec{x}' = S_2(t_1', t_2')$$

Claim For any  $f: \mathbb{F}^{k+1} \rightarrow \mathbb{F}$ ,

$$P_{\substack{S_1, S_2 \in S_k^{k+1} \\ \vec{x} \in S_1, S_2}} (\pi(S_1)(t_1, t_2) = f(\vec{x}) = \pi(S_2)(t_1', t_2')) \geq$$



$$P_{S, \vec{x} \in S} (\pi(S)(t_1, t_2) = f(\vec{x})) \quad (2)$$

Note possibly  $S_1 = S_2$ ;  
counting triplets  
 $(S_1, S_2, \vec{x})$  s.t.  
 $\vec{x} \in S_1, S_2$

$$S(t_1, t_2) = \vec{x}$$

Technique Counting + Convexity

$$|\mathbb{F}^m| \cdot \left| \left\{ S \in S_k^{k+1} \mid \exists \vec{x} \in S \right\} \right|$$

↑  
some x  
∀x this is the same

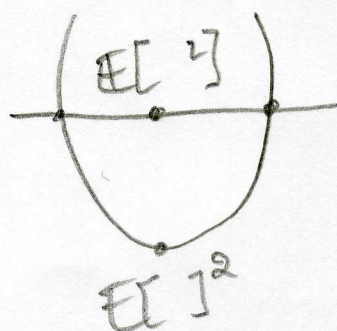
Proof Abbrer.  $I_{S,x} =$  indicator for  $\pi(S)(t_1, t_2) = f(\vec{x})$   
 where  $\vec{x} = S(t_1, t_2)$

Want to show:

$$\mathbb{E}_{\substack{S_1, S_2 \\ x \in S_1 \cap S_2}} [I_{S_1, x} \cdot I_{S_2, x}] \geq \left( \mathbb{E}_{S, x \in S} [I_{S, x}] \right)^2$$

$$\mathbb{E}_x \left[ \left( \mathbb{E}_{S \ni x} [I_{S, x}] \right)^2 \right]$$

Convexity



Inequality follows from Jensen. □

Claim For every  $S_1, S_2$  that intersect on a  $(k-1)$ -dim affine sub.,  
 $\mathbb{E}_{\vec{x} \in S_1 \cap S_2} [I_{S_1, \vec{x}} \cdot I_{S_2, \vec{x}}] > \frac{d}{|F|} \implies \pi(S_1) \& \pi(S_2)$  agree on intersec.

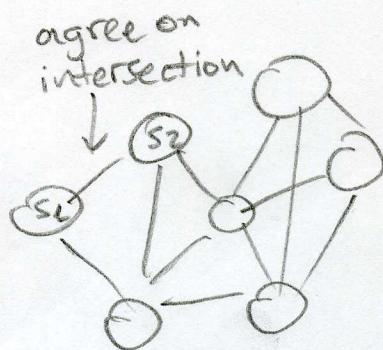
Proof Schwartz-Zippel. □

Corr  $\mathbb{P}_{\substack{S_1, S_2 \\ x \in S_1 \cap S_2}} [\pi(S_1) \& \pi(S_2) \text{ agree on intersection} + \pi(S_1) \& \pi(S_2) \text{ agree with } f \text{ on } \vec{x}] \geq \mathbb{E}_{S \in \mathcal{S}_k} \left[ \text{agree}_{\vec{x} \in S} (f, \pi(S)) \right]^2 - \frac{d}{|F|}$

Hyperplanes Graph

Vertices = hyperplanes  $\mathcal{S}_k^{(k+1)}$

Edges =  $(S_1, S_2)$  s.t.  $S_1, S_2$  agree on their intersection (possibly  $S_1 \cap S_2 = \emptyset$ ).



Note I Graph is dense by Corr.

II This is not true for dimensions  $< k$ . □