## Seminar on Sublinear Time Algorithms – Handout 2

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## 1 Today's topics

- Minimum spanning trees in sparse graphs
- Maximal and maximum matching

## 2 Open problems

• The MST approximation implies approximation within factor  $2 + \varepsilon$  for the traveling salesman problem (TSP). Can you do better?

## 3 Homework

1. Let G be a connected graph with edge weights from  $\{1, \ldots, W\}$ , and let  $G_i$  be the subgraph of G consisting of all edges with weight at most i. Prove that the cost of a minimum spanning tree can be written as

$$MST(G) = n - W + \sum_{i=1}^{W-1} c_i,$$

where  $c_i$  is the number of connected components in  $G_i$ .

Hint: Recall Kruskal's MST algorithm.

2. Let G be an undirected graph. Recall a matching is called *maximal* if it is with respect to containment, and is called *maximum* if it has the largeest possible size. Prove that every maximal matching has size which is at least half that of a maximum matching.