Take-Away Games on Beatty's Theorem and the Notion of p-Invariance

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Abstract

We formulate three reasonably short game rules for three two-pile take-away games, which share one and the same set of P-positions. This set is comprised of a pair of complementary homogeneous Beatty sequences together with (0, 0). We relate the succinctness of the game rules with the complexity of the P-positions by means of a notion dubbed p-invariance.

1 Introduction

Let us recall the rules of k-Wythoff [7], k a fixed positive integer. The available positions are (x, y), x and y non-negative integers. The legal moves are

(1A) Nim type: $(x, y) \rightarrow (x - t, y)$, if $x - t \ge 0$ and $(x, y) \rightarrow (x, y - t)$, if $y - t \ge 0$; t > 0.

(1B) Extended diagonal type: $(x, y) \rightarrow (x - s, y - t)$ if |t - s| < k and $x - s \ge 0, y - t \ge 0; s > 0, t > 0.$

This game is a so-called impartial take-away game [2], vol. 1. We restrict attention to *normal* play, that is, the player first unable to move loses. For our games it means that the player called upon to move from (0,0) loses.

Rules (1A) and (1B) imply that k-Wythoff is a so-called *invariant* [5], [11] (take-away) game, that is, each available move is legal from any position as long as the resulting position has non-negative coordinates. Every move in any invariant game is an *invariant move*.

In this note we study another type of take-away games, where certain positions have some local restrictions on the set of otherwise invariant moves. Such games are called *variant* [5], [11].

Central to our investigation is Beatty's Theorem [1] (predated by Lord Rayleigh [14]): Let $\beta > 2$ be an irrational number and define its *complement*, $\hat{\beta}$, by $\hat{\beta}^{-1} + \beta^{-1} = 1$ so that $\hat{\beta} = \beta/(\beta - 1)$. This clearly implies $1 < \hat{\beta} < 2 < \beta$. Let $A_n = \lfloor n\hat{\beta} \rfloor$, $B_n = \lfloor n\beta \rfloor$, $A = \bigcup_{n \ge 1} \{A_n\}$, $B = \bigcup_{n \ge 1} \{B_n\}$. Beatty's Theorem then asserts that A and B are *complementary* sets, that is, $A \cup B = \mathbb{Z}_{\ge 1}$, $A \cap B = \emptyset$. Since $\hat{\beta} > 1$, the (homogeneous) *Beatty* sequences (A_n) and (B_n) are strictly increasing.

1.1 Three Games

We formulate three variant games on the pairs of non-negative integers (for all games, $\beta > 2$ is irrational):

- (I) The moves are as in Nim on two piles (1A) except that the players may use an 'extended Nim type' move of the form (s,t) or (t,s) from a position, t a positive integer, $s \in \{0, \ldots, \lfloor\beta\rfloor\}$, whenever that position does not contain a coordinate in B. This game is denoted by β -Nim.
- (II) The moves are as in k-Wythoff with $k = \lfloor \beta \rfloor$ subject to (1A) and (1B), except that if one of the coordinates is in B, then only Nim-type moves (1A) are permitted. This game is denoted by β -Wynim.
- (III) The moves are as in k-Wythoff with $k = \lfloor \beta \rfloor$ subject to (1A) and (1B), except that if one of the coordinates is in B, then the move (s, t) cannot belong to the *pair of* β -triangles defined by

$$\{(x,y),(y,x) \mid (x,y) \in \{(1,\lfloor\beta\rfloor),(2,\lfloor\beta\rfloor),(2,\lfloor\beta\rfloor+1)\}\}.$$

This game is denoted by β^T -Wynim.

The main result of this note is as follows:

Theorem 1. The set of *P*-positions of β -Nim, β -Wynim and β^T -Wynim is the same. It is

$$\mathcal{P} := \bigcup_{n \ge 0} \{ (A_n, B_n) \} \bigcup_{n \ge 0} \{ (B_n, A_n) \},$$

where $A_n = \lfloor n\hat{\beta} \rfloor$, $B_n = \lfloor n\beta \rfloor$.

We prove this result in Section 3. In Section 4 we develop the distinction between invariant and variant games and relate our findings to certain complexity issues. In the section to come we give some examples.

2 Examples and tables of *P*-Positions

In the proof of the main result and in the examples of sets of P-positions to come, we use the following illustrative notation:

Notation 1. For every $n \ge 0$, (i) $\Delta A_n := A_{n+1} - A_n$, $\Delta B_n := B_{n+1} - B_n$ are the gaps. (ii) $\Delta_n := B_n - A_n$. (iii) $\Delta_n^2 := \Delta_{n+1} - \Delta_n$.

For some (invariant) take-away games on two heaps where short formulas for both the rules and the P-positions are known, such as [7, 10, 9], the coordinates of the P-positions are defined via certain algebraic numbers together with the floor function. Our first example rather uses a well known transcendental number.

Example 1. In the game of π -Wynim, a player may move as in Nim on two piles (1A), or, if the position does not contain a coordinate of the form $\lfloor \pi n \rfloor$, deviate at most $\lfloor \pi \rfloor - 1 = 2$ positions from the 'main diagonal' as given by the game k-Wythoff, that is use (1B) with k = 3. The result of this note implies that the *P*-positions of this game are the set

 $\cup_{n>0} \{ (|\hat{\pi}n|, |\pi n|), (|\pi n|, |\hat{\pi}n|) \},\$

the first few of which are displayed in Table 1.

Example 2. Example 1 illustrates Theorem 1 for a member of our second game family, β -Wynim. A further example: Let k be a positive integer, $\beta = (2 + k + \sqrt{k^2 + 4})/2$, $\hat{\beta} = \beta - k$. Then $\beta = \sqrt{2} + 2$, $\hat{\beta} = \beta - 2$; note that $\lfloor \beta \rfloor = 3$ as in Example 1. The first few *P*-positions for k = 2 are shown in Table 2. Since $\beta - \hat{\beta} = k = 2$, we have $\Delta_n = kn = 2n$, so $\Delta_n^2 = k = 2$ for all $n \ge 0$.

Table 1: The first few *P*-positions (A_n, B_n) for β -Nim, β -Wynim and β^T -Wynim, $\beta = \pi = 3.14159...$

Remark 1. It is remarkable that, for $\beta = (2 + k + \sqrt{k^2 + 4})/2$, k-Wythoff has the same set of *P*-positions as our three games. In particular, for k = 1, 1-Wythoff is the classical Wythoff game [2]. For k = 2, the first few *P*positions of both games are displayed in Table 2. In [4] it was shown that from the classical Wythoff game no move can be deleted while preserving the set of *P*-positions of the classical Wythoff game. In the present note, Wythoff moves were deleted, and the *P*-positions are still preserved. The difference is that in [4] only *invariant* moves were permitted. See section 4 for more on the latter topic.

3 Proof of the Main Result

We preface the proof of Theorem 1 by collecting some facts on the sets $\{A_n\}$ and $\{B_n\}$.

Proposition 1. For every $n \ge 0$, (i) The only possible gap pairs are

 $(\Delta A_n, \Delta B_n) \in \{(1, \lfloor \beta \rfloor), (1, \lfloor \beta \rfloor + 1), (2, \lfloor \beta \rfloor), (2, \lfloor \beta \rfloor + 1)\}.$

(ii) $\Delta_n^2 = \Delta B_n - \Delta A_n.$ (iii) $\Delta_n^2 \in \{\lfloor \beta \rfloor - 2, \lfloor \beta \rfloor - 1, \lfloor \beta \rfloor\}.$ Table 2: The first few *P*-positions (A_n, B_n) for 2-Wythoff, β -Nim β -Wynim and β^T -Wynim; β as in Example 2.

Proof. (i) This is a well known result.

(ii) $\Delta_n^2 = (B_{n+1} - A_{n+1}) - (B_n - A_n) = (B_{n+1} - B_n) - (A_{n+1} - A_n) = \Delta B_n - \Delta A_n.$

(iii) Follows directly from (i) and (ii).

Example 3. Notice that in Example 1, Table 1, Δ_n^2 assumes all three possible values $\{1, 2, 3\} = \{\lfloor \beta \rfloor - 2, \lfloor \beta \rfloor - 1, \lfloor \beta \rfloor\}$. In Example 2, Δ_n^2 assumes only the value $2 = \lfloor \beta \rfloor - 1$.

Proof of Theorem 1. Since our games are acyclic, it suffices to demonstrate the following two properties for each game:

 $P \rightarrow N$: Every move from any position of the form

$$(A_n, B_n) \text{ or } (B_n, A_n) \tag{1}$$

results in a position outside (1).

 $N \to P$: Given any position outside (1), there exists a move into (1).

For the direction $P \to N$ we use the same argument for the games (I) β -Nim and (II) β -Wynim, namely: Suppose that we play from a position of the form (1). The game rules imply that only Nim type moves (1A) are permitted so that by complementarity, there is no move to a position of the same form.

For game (III) β^T -Wynim, we have to show that both (i) $(A_n, B_n) \rightarrow (A_m, B_m)$ and (ii) cross moves $(A_n, B_n) \rightarrow (B_m, A_m)$ are blocked for every $0 \leq m < n$.

(i) By Proposition 1, $(B_n - B_m) - (A_n - A_m) = \Delta_n - \Delta_m \ge \Delta_n - \Delta_{n-1} = \Delta_{n-1}^2 \in \{\lfloor \beta \rfloor - 2, \lfloor \beta \rfloor - 1, \lfloor \beta \rfloor\}$ where the \ge follows from the fact that $\beta > \hat{\beta}$, which implies that Δ_i is a non-decreasing function of *i*. Therefore the move $(A_n, B_n) \to (A_m, B_m)$ is either blocked by the triangle move restriction of β^T -Wynim (if $\Delta_{n-1}^2 \le \lfloor \beta \rfloor$), or by the $\lfloor \beta \rfloor$ -Wythoff constraint (if $\Delta_{n-1}^2 \ge \lfloor \beta \rfloor$).

(ii) Notice that this move is possible only if $A_n > B_m$. Now $(B_n - A_m) - (A_n - B_m) = \Delta_n + \Delta_m$. Similarly to (i), if $\Delta_n + \Delta_m \in \{\lfloor \beta \rfloor - 2, \lfloor \beta \rfloor - 1, \lfloor \beta \rfloor\}$, this forces m = 0 and n = 1 so that the move is blocked by the β -triangle move restriction; otherwise by the $|\beta|$ -Wythoff constraint.

For the direction $N \to P$, let (x, y), $0 \le x \le y$ be a position not of the form (1). We assume first that this position has a coordinate of the form B_n , so for each game it suffices to show that a Nim type (1A) move suffices for moving into (1). If $x = B_n$ then move $y \to A_n$. If $y = B_n$ and $x > A_n$ then move $x \to A_n$. If $y = B_n$ and $x < A_n$, complementarity implies that there exists m < n such that either $x = A_m$ so the move $y \to B_m$ restores (1); or else $x = B_m$, so the move $y \to A_m$ does it.

Hence we may assume that both x and y are in A, say $x = A_m \leq A_n = y$. If $y > B_m$, then the Nim type (1A) move $y \to B_m$ suffices for each game. We may therefore assume that

$$x = A_m \le A_n = y < B_m. \tag{2}$$

Since each of (A_i) and (B_i) is strictly increasing, a Nim type move to a position (1) does not exist, so we have to find a (1B) extended diagonal type move for the games β -Wynim and β^T -Wynim. Observe that for both these games, this type of moves is now unrestricted with $k = |\beta|$.

Let d := y - x. Then $d = A_n - A_m < B_m - A_m = \Delta_m$. By Proposition 1, Δ_i grows from 0 to Δ_m as *i* grows from 0 to *m*, in steps $\Delta_i^2 = \Delta B_i - \Delta A_i \in \{\lfloor \beta \rfloor - 2, \lfloor \beta \rfloor - 1, \lfloor \beta \rfloor\}$, bounded above by $\lfloor \beta \rfloor$. Hence there exists *j* such that $0 \leq d - \Delta_j < \lfloor \beta \rfloor$. Then move $(x, y) \to (A_j, B_j)$. We need to show three things: (i) j < m, (ii) $y > B_j$, (iii) $|(y - B_j) - (A_m - A_j)| < \lfloor \beta \rfloor$.

(i) $\Delta_j \leq d = y - A_m < B_m - A_m = \Delta_m$. Since Δ_i is an increasing function

of i, we have j < m.

(ii)
$$y = A_m + d > A_j + d \ge A_j + \Delta_j = B_j.$$

(iii)
$$|(y - B_j) - (A_m - A_j)| = |(y - A_m) - (B_j - A_j)| = |d - \Delta_j| < \lfloor \beta \rfloor.$$

On the other hand, for the game β -Nim and a position of the form in (2), by Proposition 1 (i) a nearest lower P-position is attainable by an extended 'horizontal' Nim-type move. Precisely, since $\Delta B_n \in \{\lfloor \beta \rfloor, \lfloor \beta \rfloor + 1\}$ we can lower $y = A_n$ to B_i , where $i \ge 0$ is such that $B_i < A_n < B_{i+1}$, and $x = A_m$ to A_i , that is move $(A_m, A_n) \to (A_i, B_i)$. By the (1A) Nim type move we have to show that $A_i < A_m$. But the definition of i together with (2) give $B_i < A_n < B_m$ which, by i < m, implies $A_i < A_m$.

Thus the set \mathcal{P} is indeed the set of P-positions for our three games.

4 *p*-Invariance and game complexity

Let us continue our brief discussion of variant versus invariant games from the introduction and Remark 1 and relate it to the complexity of numbers and games. The complexity of a real number is its status as an integer, rational, algebraic or transcendental number.

Our three games are, in fact, 'minimally variant' in the sense that all their positions can be partitioned into precisely two sets, namely,

$$\{(A_i, A_i)\}$$
 and $\{(x, y) \mid x = B_i \text{ or } y = B_i \text{ for some } i\},\$

such that, for each game, for each set, the possible moves are invariant. This observation motivates a weakening of the notion of invariance to *p*-invariance, $p \in \mathbb{Z}_{>0}$, for certain variant games. A game is *p*-invariant if its positions can be partitioned into *p* subsets, such that every move from any position in each subset is invariant, and it cannot be thus partitioned into less than *p* subsets. An invariant game is 1-invariant.

Let $\gamma = \beta - \hat{\beta}$. It appears that the complexity of γ , the simplicity of the game rules and the size of p are related. By complexity of a real number, we mean its category: integer, rational, algebraic or transcendental. If $\gamma = k$ is an integer, there are simple game rules, (1A) and (1B), and the game is invariant (Example 2). For our three games, γ is not necessarily an integer, the game rules are longer and the 1-invariance is replaced by 2-invariance. To shed more light on these suggested relationships, it might be well to

investigate whether γ rational, algebraic [5, 12], or transcendental has any effect on the length of the game rules and *p*-invariance.

We close this section with a conjecture which requires a little background.

The succinct input size of a given ordered pair of integers (x, y) is $\log(xy)$. The time complexity of deciding whether a given ordered pair (x, y) is of the form (A_n, B_n) is polynomial in $\log(xy)$, see [7], §3. In [11, Main Theorem] it is demonstrated that, given the set \mathcal{P} in Theorem 1, there is an *invariant* game for which the time complexity of determining whether a given ordered pair (s, t) is a legal move is exponential in $\log(st)$. In [5, 12] polynomial time invariant game rules are determined for the set \mathcal{P} when γ is restricted to some specific algebraic numbers of degree 2.

We make the following conjecture.

Conjecture 1. Given a set $\mathcal{P} = \bigcup_{n\geq 0} \{(A_n, B_n)\} \bigcup_{n\geq 0} \{(B_n, A_n)\}$ where $(A_n)_{n\geq 1}$ and $(B_n)_{n\geq 1}$ are complementary Beatty sequences with $\beta > 2$ transcendental, and a position $(x, y) \neq (0, 0)$, $(x, y) \notin \mathcal{P}$. Then there is no invariant game with polynomial time complexity in $\log(xy)$ for finding a move from (x, y) into a *P*-position.

Perhaps this conjecture holds even if γ is algebraic of degree > 2, many γ of degree 2, perhaps even if γ is a non-integer rational number.

The notion of p-invariance is also interesting in a somewhat different context. In [13] certain p-invariant 2-heap subtraction games with a finite number of moves are studied and it is showed that they embrace computational universality.

5 Discussion

We have formulated three reasonably short game rules for three 2-invariant games, which have identical sets of *P*-positions. Suppose that we fix β and then increase the density of the pairs of sequences from 1 to say an arbitrary number $\gamma > 1$ (or decreases to a density < 1) where α is defined via $1/\alpha + 1/\beta = \gamma$. (That is, for all β , $\alpha \neq \hat{\beta}$.) Given candidate *P*-positions as above, is there still a short/succinct but non-trivial way of formulating the game rules without disclosing both irrationals or/and the joint density of the sequences? It is unknown to us whether or not there exist invariant rules for such games, see [11], [6]. Is it possible to find 2-invariant rules in the sense of this note? As another remark, observe that neither $\hat{\beta}$ nor the density 1 is disclosed in the presentation of the rules of our games in this note. In [11] invariant game rules are given for candidate *P*-positions constructed from complementary Beatty sequences, but not in a single case have we found a succinct description. In this note we have chosen to remove the nice condition of invariance (and reverted to 2-invariance) from the game rules and, maybe even more notably, one of the coordinates of the candidate P-positions is disclosed within the game rules. This could be argued to be a severe drawback in a definition of the rules of a game. But, on the other hand, we were able to give a very succinct formulation, without a complete trivialization of game rules, for all complementary homogeneous Beatty sequences and these are uncountably many.

For other investigations that relate Nim and Wythoff, see [3], [8].

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