

## Seminar on Algorithms and Geometry

### Lecture 5

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## 1 Johnson-Lindenstrauss Lemma and Concentration of Measure

In the previous lecture we stated and proved the following theorem

**Theorem 1 (Johnson-Lindenstrauss)** *For every subset  $X \subseteq \ell_2$  and every  $\epsilon \geq 0$  there is an embedding  $f : X \hookrightarrow \ell_2^k$  with distortion  $1 + \epsilon$  and dimension  $k = O(\frac{1}{\epsilon^2} \log n)$ .*

In this lecture we will see a sketch of an alternative proof of the theorem with an emphasis on the phenomenon of concentration of measure.

**Theorem 2** *Let  $L$  be a random subspace of  $\mathbb{R}^n$  of dimension  $k$  and let  $f : \mathbb{R}^n \mapsto \mathbb{L}$  be an orthogonal projection onto  $L$  (here we think of  $L$  as a copy of  $\mathbb{R}^k$ ). Then there exists a constant  $c = c(n, k)$  s.t. for every  $x, y \in \mathbb{R}^n$*

$$\Pr_f[1 - \epsilon \leq \frac{\|f(x) - f(y)\|}{c\|x - y\|} \leq 1 + \epsilon] \geq 1 - \frac{1}{n^3}$$

**Sketch of Proof**  $L$  is chosen by picking  $k$  orthogonal vectors from  $\mathbb{S}^{n-1}$  s.t. each vector has a uniform distribution over  $\mathbb{S}^{n-1}$ . Alternatively choose a random rotation  $U : \mathbb{R}^n \mapsto \mathbb{R}^n$  and let  $L = U(\text{span}\{e_1, \dots, e_k\})$  where  $e_1, \dots, e_k$  are the first  $k$  vectors of the canonical basis of  $\mathbb{R}^n$ . As in the previous class, it suffices to prove that for all  $v \in \mathbb{S}^{n-1}$

$$\Pr_f[1 - \epsilon \leq \frac{\|f(v)\|}{c} \leq 1 + \epsilon] \geq 1 - \frac{1}{n^3}$$

Denote  $w = (w_1, \dots, w_n) = U^{-1}v$  and note that the projection  $v$  onto  $L$  has the same length as that of  $w$  onto  $U^{-1}(L) = \text{span}\{e_1, \dots, e_k\}$  so

$$\|f(v)\| = \|(w_1, \dots, w_k)\| = \sqrt{w_1^2 + \dots + w_k^2}$$

Let  $\mu$  denote the uniform probability measure on  $\mathbb{S}^{n-1}$  (the Haar measure), then  $w = U^{-1}v$  is distributed according to  $\mu$ . We now ask what is the length of the projection of  $w$  on the first  $k$  coordinates.

**Theorem 3** *Let  $g : \mathbb{S}^{n-1} \mapsto \mathbb{R}$  be 1-Lipschitz i.e.*

$$|g(x) - g(y)| \leq \|x - y\|$$

and let  $m = m(g)$  be a median of  $g$  i.e.

$$\Pr_{x \in \mu}[g(x) \geq m] \geq \frac{1}{2} \text{ and } \Pr_{x \in \mu}[g(x) \leq m] \geq \frac{1}{2}$$

then for all  $\delta > 0$

$$\Pr_{x \in \mu}[|g(x) - m| \geq \delta] \leq 4e^{-\delta^2 n/2}$$

We can use theorem 3 to prove theorem 2. It can be easily verified that  $g(w) = \|(w_1, \dots, w_k)\|$  is 1-Lipschitz. Applying theorem 3 gives us the following bound

$$\Pr_w[m - \delta \leq g(w) \leq m + \delta] \geq 1 - 4e^{-\delta^2 n/2}$$

We now choose  $c = c(n, k) = m$  and  $\delta = \epsilon m$  to get

$$\begin{aligned} \Pr_f[1 - \epsilon \leq \frac{\|f(v)\|}{c} \leq 1 + \epsilon] = \\ \Pr_w[m(1 - \epsilon) \leq g(w) \leq m(1 + \epsilon)] \geq 1 - 4e^{-\epsilon^2 m^2 n/2} \end{aligned}$$

All that is left is to lower bound  $m$ . We observe that

$$\mathbb{E}_{w \in \mu}[g(w)^2] = \mathbb{E}[w_1^2 + \dots + w_k^2] = k\mathbb{E}[w_1^2] = \frac{k}{n}\mathbb{E}[w_1^2 + \dots + w_k^2] = \frac{k}{n}$$

where we have used the symmetry of the coordinates, all of the coordinates are identically distributed. We can use this to lower bound the median. Consider a parameter  $t > 0$  we partition the integration into two parts  $[0, (m+t)]$  and  $[(m+t), 1]$  and upper bound each part

$$\begin{aligned} \frac{k}{n} &= \mathbb{E}[g(w)^2] \\ &\leq \Pr[g(w)^2 \geq (m+t)^2] \cdot 1 + \Pr[g(w)^2 \leq (m+t)^2] \cdot (m+t)^2 \\ &\leq 4e^{-t^2 n/2} + (m+t)^2 \end{aligned}$$

By choosing  $t = \sqrt{\frac{k}{5n}}$  we get  $m \geq \Omega(\sqrt{\frac{k}{n}})$ . Finally

$$\Pr_f[1 - \epsilon \leq \frac{\|f(v)\|}{c} \leq 1 + \epsilon] \geq 1 - e^{-\epsilon^2 k c'} \geq 1 - \frac{1}{n^3}$$

if  $k \geq 100 \frac{1}{\epsilon^2} \log n$ . ■

How can we prove theorem 3? The following isoperimetric inequalities provide an answer.

**Theorem 4 (Paul Levy 1951)** *Let  $A \subseteq \mathbb{S}^{n-1}$  be a measurable set and let  $B \subseteq \mathbb{S}^{n-1}$  be a cap with  $\mu(A) = \mu(B)$ . Then for all  $\epsilon > 0$ ,  $\mu(A_\epsilon) \geq \mu(B_\epsilon)$ . Here*

$$A_\epsilon = \{x \in \mathbb{S}^{n-1} : d(x, A) \leq \epsilon\}$$

and similarly for  $B_\epsilon$ .

Remark: Here distance is Euclidean, i.e. measured according to  $\ell_2$ -norm, but similar theorems can be proved for geodesic distance on the sphere.

Using Theorem 4 plus estimates on the volume of a spherical cap, one can obtain the following bound on the measure of  $A_\epsilon$ . Such a bound can also be proved directly via the Brunn-Minkowski Theorem.

**Theorem 5** *Let  $A \subseteq \mathbb{S}^{n-1}$  be a measurable set with  $\mu(A) \geq \frac{1}{2}$  then for all  $\epsilon > 0$*

$$\mu(A_\epsilon) \geq 1 - 2e^{-\epsilon^2 n/2}$$

Consider cutting the sphere by a hyperplane passing through the origin. The sets on both sides have a measure of exactly  $\frac{1}{2}$ . Applying the inequality to each one of these sets we see that almost all of the measure is concentrated on a thin strip around the equator. The total amount of measure outside is an exponentially small function of the dimension.

We now sketch the proof of theorem 3 using theorem 5.

**Sketch of Proof** Apply theorem 5 to  $A^- = \{x \in \mathbb{S}^{n-1} : g(x) \leq m\}$  and obtain a lower bound on  $\mu(A_\epsilon^-)$ . Then do similarly for  $A^+ = \{x \in \mathbb{S}^{n-1} : g(x) \geq m\}$ . ■