

# Advanced Algorithms 2012A – Problem Set 7

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## Extra credit:

1. Design a variant of Cheeger's inequalities, where the input graph comes with two distinguished vertices  $s, t \in V$ , and we only consider cuts that separate  $s$  from  $t$ . Specifically, find a value  $\lambda_{st}(G)$  that satisfies something like:

$$\frac{1}{2}\lambda_{st}(G) \leq \min_{S: |\{s,t\} \cap S|=1} \phi_G(S) \leq \sqrt{2\lambda_{st}(G)}$$

Remark: This  $\lambda_{st}(G)$  need not be an eigenvalue, although it is derived from the (normalized) Laplacian matrix.

If helpful, you may assume that  $G$  is  $d$ -regular.

2. Let  $G = (V, E)$  be a capacitated graph (i.e., with edge capacities  $c_e \geq 0$ ). Show that the algorithm below outputs a capacitated tree  $T$  on the same vertex set  $V$  that is *flow-equivalent* to  $G$ , in the sense that for all  $s, t \in V$ , the maximum  $st$ -flow in  $G$  is equal to that in  $T$ . (Notice that  $T$  has same vertices, but in general different edges and capacities than  $G$ .)

Algorithm:

- (1) Construct on  $V$  a complete graph  $G'$ , with every edge capacity  $c'_{st}$  is equal to  $s, t \in V$  are connected by an edge whose capacity  $c'_{st}$  is defined to be the maximum  $st$ -flow in  $G$ .
- (2) Compute a maximum spanning tree  $T$  for  $G'$  (with respect to the edge capacities  $c'_e$ ).

Hint: A maximum spanning tree is analogous to a minimum spanning tree; in particular, it is just a minimum spanning tree for the capacities  $-c'_e$ .