

# Lecture 8: Cosystolic expansion: local to global

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In this lecture we will prove cosystolic expansion for bounded-degree complexes, via a local to global connection. We will show that if a simplicial complex has links that are coboundary expanders, together with spectral expansion, this implies global cosystolic expansion. This theorem was proven by Kaufman Kazhdan and Lubotzky [3] for dimension 1, and by Evra and Kaufman [1] for higher dimensions.

Cosystolic expansion, as we have mentioned before, is related to

- Property testing (we have seen, for example, that coboundary expansion of the complete complex is the same as bi-clique testability)
- Locally testable codes and PCPs
- Topological overlap property (TOP)
- Stability of covers

## 1 Definitions

Recall from lecture 4 that the  $i$ -coboundaries are  $B^i = \text{Im}(\delta_{i-1})$  and the  $i$ -cocycles are  $Z^i = \text{Ker}(\delta_i)$ . We noted that  $B^i \subseteq Z^i \subseteq C^i$  and defined the  $i$ -th cohomology to be  $H^i = Z^i/B^i$ .

A simplicial complex  $X$  is a  $\beta$ -cosystolic expander in dimension  $i$  if for every  $i$ -chain  $f \in C^i(X, \mathbb{Z}_2)$ ,

$$\text{wt}(\delta_i f) \geq \beta \cdot \text{dist}(f, B^i).$$

If, furthermore,  $H^i = 0$ , we say that  $X$  is a  $\beta$ -coboundary expander in dimension  $i$ . We also denote by  $h^i$  the largest  $\beta$  for which  $X$  is a  $\beta$ -cosystolic expander:

$$h^i(X, \mathbb{Z}_2) = \min_{f \in C^i \setminus B^i} \frac{\text{wt}(\delta_i f)}{\text{dist}(f, B^i)}.$$

Is there a “Cheeger’s inequality” for high dimensional expanders? One could hope that every spectral HDX is a cosystolic expander, and possibly also vice versa. However, this is not true. Gundert and Wagner [2] give a random construction of a complex which is a very good spectral expander yet the cosystolic constant is  $o(1)$ .

## 2 Local to Global: cosystolic expansion comes from expansion in the links

**Theorem 2.1.** *Let  $X$  be a 3 dimensional simplicial complex. Suppose that*

- $X$  is a  $\gamma$ -two-sided link expander.
- For every  $v \in X(0)$ ,  $X_v$  is a  $\beta_0$ -coboundary expander in dimension 1.

Then,  $X$  is a  $\beta$  cosystolic expander in dimension 1 for some constant  $\beta$ .

Observe that the statement makes sense only if the dimension of the link is at least 2 (otherwise we cannot speak of coboundary expansion of 1-chains), so the complex needs to be 3 dimensional for such a statement to hold. Are there 2-dimensional HDX that are not 2-skeletons of 3 dimensional HDX and yet are cosystolic expanders? This is not known.

The theorem can be generalized to expansion in any dimension  $i$ , as long as the complex is  $i + 2$  dimensional [1].

*Proof.* Let  $f \in C^1(X, \mathbb{Z}_2)$ , and let  $\varepsilon = wt(\delta f)$ . We will show that there is some  $\tilde{f} \in Z^1$  such that

$$\beta' \cdot \text{dist}(f, \tilde{f}) \leq wt(\delta f).$$

We use the following local correction algorithm:

**Algorithm.** If there is a vertex such that changing  $f$  on all edges touching  $v$  can reduce  $wt(\delta f)$ , change it. Repeat

Observe that the algorithm must halt because every iteration reduces the size of  $\delta f$  by at least one. Let  $f$  be the initial chain, and let  $\tilde{f}$  be the chain after the algorithm halts.

**Claim 2.2.** There is some constant  $d_0$  such that  $\text{dist}(f, \tilde{f}) < \varepsilon \cdot d_0$ .

*Proof.* The number of iterations of the algorithm is at most the support of  $\delta f$ , which we denote by  $|\delta f| = \varepsilon |X(2)|$ . So the number of edges that are modified during the course of the algorithm is at most  $|\delta f| = \varepsilon |X(2)| = \varepsilon |X(1)| \cdot d_0$ , where  $d_0 = \frac{|X(2)|}{|X(1)|}$ .  $\square$

Our main lemma is the following

**Lemma 2.3.** If  $\delta \tilde{f} \neq 0$  then  $wt(\delta \tilde{f}) > \tau$ .

This implies the theorem with  $\beta = \min(\tau, \frac{1}{d_0})$  because if  $\delta \tilde{f} = 0$  then the above claim gives the required inequality, and if  $\delta \tilde{f} \neq 0$  then  $wt(\delta f) > \tau \cdot 1 \geq \tau \text{dist}(f, \tilde{f})$ .  $\square$

## 2.1 Proof of Lemma 2.3

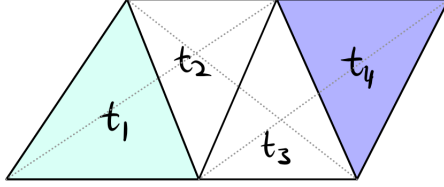
In this section let  $f$  denote the chain after the algorithm ended. Let

$$T^* = \{t \in X(2) \mid \delta f(t) \neq 0\}$$

The heart of the matter is to look at redundancies between the triangle constraints, and use them to propagate errors. Consider a 3-face  $p = \{a, b, c, d\}$ . It contains four triangles and six edges. Moreover, no matter when the value of  $f$  on the edges is, the number of triangles that belong to  $T^*$  must be even. So if it is non-zero, it must be larger than one.

Now consider an upper random walk

$$t_1 \rightarrow t_2 \rightarrow t_3 \rightarrow t_4$$



where every consecutive pair of triangles belong together to a 3-face. (This implies that they intersect on an edge). If we know that  $t_1 \in T^*$ , there is a probability of at least  $1/3$  that  $t_2 \in T^*$ , because in the 3-face  $t_1 \cup t_2$  there are at least two triangles for which  $\delta f \neq 0$ . For the same reason, there is probability  $\frac{1}{3} \cdot \frac{1}{3} = \frac{1}{9}$  that  $t_3 \in T^*$  and probability  $\frac{1}{3} \cdot \frac{1}{3} \cdot \frac{1}{3} = \frac{1}{27}$  that  $t_4 \in T^*$ .

If this were the case *even conditioned on*  $t_1 \cap t_4 = \emptyset$ , then we would get, via swap-walk arguments, that the set  $T^*$  is large.

**Claim 2.4.** If  $\mathbb{P}_{t_1, t_2, t_3, t_4} [t_4 \in T^* | t_4 \cap t_1 = \emptyset] \geq \mathbb{P}[t_1 \in T^*] \cdot \alpha$  then  $|T^*| = \Omega(X(2))$ .

The reason is that the graph whose nodes are the triangles, and where we connect two triangles  $t_1, t_4$  according to the rule above, is an expander graph. So by applying the Alon-Chung lemma we get that a subset of nodes with average degree above  $\alpha$  must be linearly large as long as  $\alpha > \lambda$ .

However, the proof is more complicated, because there is a constant fraction of random walks  $t_1 \rightarrow t_2 \rightarrow t_3 \rightarrow t_4$  in which all of  $t_1, t_2, t_3, t_4$  share an edge, or a vertex. If  $t_1 \in T^*$  it could a priori be that all of the walks in which also  $t_4 \in T^*$  are such walks. In that case the above reasoning would fail. However, when this happens it means that either

- One of the edges touching  $t_1$  is heavy (touches many triangles in  $T^*$ ), or
- One of the vertices touching  $t_1$  is heavy (touches many triangles in  $T^*$ )

So to complete the proof we show that the set of heavy vertices and edges is tiny, even with respect to  $T^*$  and therefore, most of the time the walks don't get stuck.

Why are there few heavy vertices, for example? For this we must look at the triangles in its link

$$T_v^* = \{t \in T^* \mid t \cup v \in X(3)\}.$$

**Claim 2.5.** If  $T_v^*$  is small, we can use coboundary expansion of  $X_v$  to find a better assignment to the edges touching  $v$ .

In addition, we can show, using the fact that the bipartite graph connecting a vertex  $v$  to a triangle  $t \in T_v$  is a very strong expander, that only very few vertices  $v$  can have large  $T_v^*$  (since all in all the set  $T^*$  is by assumption not too large). So only those vertices can be heavy vertices, and these make up a tiny fraction of all vertices.

A similar, but more subtle, argument can be made for edges. (We show that except for a tiny minority, a heavy edge almost always touches a heavy vertex).

### 3 Is the extra dimension really needed?

One might conjecture that link-expansion suffices for proving coboundary expansion. This would give a "Cheeger inequality" in dimensions higher than 0. However, Gundert and Wagner [2] give a counterexample. They describe a randomized construction of a 2-dimensional simplicial complex whose links are excellent spectral expanders; and yet the complex has poor coboundary expansion.

### References

- [1] Shai Evra and Tali Kaufman. Bounded degree cosystolic expanders of every dimension. In *Proceedings of the 48th Annual ACM SIGACT Symposium on Theory of Computing, STOC 2016, Cambridge, MA, USA, June 18-21, 2016*, pages 36–48, 2016. [1](#), [2](#)
- [2] Anna Gundert and Uli Wagner. On eigenvalues of random complexes. *Israel Journal of Mathematics*, 216(2):545–582, 2016. [1](#), [4](#)
- [3] Tali Kaufman, David Kazhdan, and Alexander Lubotzky. Ramanujan complexes and bounded degree topological expanders. In *Proc. 55th IEEE Symp. on Foundations of Comp. Science (FOCS)*, pages 484–493, 2014. [1](#)