### Survey: Compressive Sensing in Signal Processing

Justin Romberg

Georgia Tech, School of ECE

Sublinear Algorithms

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# Acquisition as linear algebra



- Small number of samples = underdetermined system Impossible to solve in general
- If x is *sparse* and  $\Phi$  is *diverse*, then these systems can be "inverted"

DSP: sample first, ask questions later

Explosion in sensor technology/ubiquity has caused two trends:

- Physical capabilities of hardware are being stressed, increasing speed/resolution becoming *expensive* 
  - gigahertz+ analog-to-digital conversion
  - accelerated MRI
  - industrial imaging
- Deluge of data
  - camera arrays and networks, multi-view target databases, streaming video...

Compressive Sensing: sample smarter, not faster

# Sparsity/Compressibility

 $N \\ {\rm pixels}$ 



 $S \ll N$ large wavelet coefficients

N wideband signal samples



 $S \ll N$ large Gabor coefficients

time

# Wavelet approximation



1 megapixel image



25k term approximation

• 1% error with  $\approx 2.5\%$  of the wavelet coefficients



• If x is *sparse* and  $\Phi$  is *diverse*, then these systems can be "inverted"

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 $y = Ax_0 + \text{noise}$ 

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• Standard way to recover  $x_0$ , use the *pseudo-inverse* 

solve 
$$\min_{x} \|y - Ax\|_2^2 \quad \Leftrightarrow \quad \hat{x} = (A^T A)^{-1} A^T y$$

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• A: When the matrix A is an *approximate isometry*...

$$||Ax||_2^2 \approx ||x||_2^2$$
 for all  $x \in \mathbb{R}^N$ 

i.e. A preserves *lengths* 

 Suppose we have an M × N observation matrix A with M ≥ N (MORE observations than unknowns), through which we observe

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$$||A(x_1 - x_2)||_2^2 \approx ||x_1 - x_2||_2^2$$
 for all  $x_1, x_2 \in \mathbb{R}^N$ 

i.e. A preserves *distances* 

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• A: When the matrix A is an *approximate isometry*...

$$(1-\delta) \le \sigma_{\min}^2(A) \le \sigma_{\max}^2(A) \le (1+\delta)$$

i.e. A has clustered singular values

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$$(1-\delta)\|x\|_2^2 \leq \|Ax\|_2^2 \leq (1+\delta)\|x\|_2^2$$

for some  $0 < \delta < 1$ 

• Now we have an underdetermined  $M \times N$  system  $\Phi$  (FEWER measurements than unknowns), and observe

 $y = \Phi x_0 + \text{noise}$ 

When can we stably recover an S-sparse vector?

• Now we have an underdetermined  $M \times N$  system  $\Phi$  (FEWER measurements than unknowns), and observe

$$y = \Phi x_0 + \text{noise}$$

• We can recover  $x_0$  when  $\Phi$  is a *keeps sparse signals separated* 

$$(1-\delta) \|x_1 - x_2\|_2^2 \leq \|\Phi(x_1 - x_2)\|_2^2 \leq (1+\delta) \|x_1 - x_2\|_2^2$$

for all S-sparse  $x_1, x_2$ 

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• To recover  $x_0$ , we solve

 $\min_{x} \ \|x\|_0 \quad \text{subject to} \quad \Phi x \approx y$ 

 $||x||_0 =$  number of nonzero terms in x

• This program is intractable

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• A relaxed (convex) program

 $\min_{x} ||x||_1 \text{ subject to } \Phi x \approx y$ 

 $||x||_1 = \sum_k |x_k|$ 

• This program is very tractable (linear program)

# Sparse recovery algorithms

- Given y, look for a sparse signal which is consistent.
- One method:  $\ell_1$  minimization (or *Basis Pursuit*)

$$\min_{x} \|\Psi^T x\|_1 \quad \text{s.t.} \quad \Phi x = y$$

 $\Psi =$ sparsifying transform,  $\Phi =$  measurement system (need RIP for  $\Phi\Psi$ )

Convex (linear) program, can relax for robustness to noise

Performance has theoretical guarantees

• Other recovery methods include greedy algorithms and iterative thresholding schemes

# Stable recovery

- Despite its nonlinearity, sparse recovery is stable in the presence of
  - modeling mismatch (approximate sparsity), and
  - measurement error
- If we observe  $y = \Phi x_0 + e$ , with  $||e||_2 \le \epsilon$ , the solution  $\hat{x}$  to

$$\min_{x} \|\Psi^T x\|_1 \quad \text{s.t.} \quad \|y - \Phi x\|_2 \le \epsilon$$

will satisfy

$$\|\hat{x} - x_0\|_2 \leq \operatorname{Const} \cdot \left(\epsilon + \frac{\|x_0 - x_{0,S}\|_1}{\sqrt{S}}\right)$$

where

- $x_{0,S} = S$ -term approximation of  $x_0$
- $\blacktriangleright~S$  is the largest value for which  $\Phi\Psi$  satisfies the RIP
- Similar guarantees exist for other recovery algorithms
  - greedy (Needell and Tropp '08)
  - iterative thresholding
     (Blumensath and Davies '08)

• They are very hard to design, but they exist everywhere!



#### • For any fixed $x \in \mathbb{R}^N$ , each measurement is

 $y_k \sim \operatorname{Normal}(0, \|x\|_2^2/M)$ 

• They are very hard to design, but they exist everywhere!



• For any fixed  $x \in \mathbb{R}^N$ , we have

$$\mathbf{E}[\|\Phi x\|_2^2] = \|x\|_2^2$$

the mean of the measurement energy is exactly  $||x||_2^2$ 

• They are very hard to design, but they exist everywhere!



• For any fixed  $x \in \mathbb{R}^N$ , we have

$$P\left\{\left|\|\Phi x\|_{2}^{2}-\|x\|_{2}^{2}\right| < \delta \|x\|_{2}^{2}\right\} \geq 1-e^{-M\delta^{2}/4}$$

• They are very hard to design, but they exist everywhere!



• For all 2S-sparse 
$$x \in \mathbb{R}^N$$
, we have  

$$P\left\{\max_x \left|\|\Phi x\|_2^2 - \|x\|_2^2\right| < \delta \|x\|_2^2\right\} \ge 1 - e^{c \cdot S \log(N/S)} e^{-M\delta^2/4}$$
So we can make this probability close to 1 by taking  
 $M \gtrsim S \log(N/S)$ 

Four general frameworks:

- Random matrices (iid entries)
- Random subsampling
- Random convolution
- (Randomly modulated integration we'll skip this today)

Note the role of randomness in all of these approaches

Slogan: random projections keep sparse signal separated

# Random matrices (iid entries)



- Random matrices are provably efficient
- We can recover S-sparse x from

$$M \gtrsim S \cdot \log(N/S)$$

measurements

### Rice single pixel camera



(Duarte, Davenport, Takhar, Laska, Sun, Kelly, Baraniuk '08)

# Georgia Tech analog imager



### Compressive sensing acquisition







### Random matrices

Example:  $\Phi$  consists of *random rows* from an *orthobasis* U



Can recover S-sparse x from

$$M \gtrsim \mu^2 S \cdot \log^4 N$$

measurements, where

$$\mu = \sqrt{N} \max_{i,j} |(U^T \Psi)_{ij}|$$

is the *coherence* 

# Examples of incoherence

• Signal is sparse in time domain, sampled in Fourier domain



S nonzero components

freq domain  $\hat{x}(\omega)$ 

measure m samples

• Signal is sparse in wavelet domain, measured with noiselets

example noiselet

wavelet domain



(Coifman et al '01)

noiselet domain



### Accelerated MRI



(Lustig et al. '08)

Empirical processes and structured random matrices

• For matrices with this type of *structured randomness*, we simply do not have enough concentration to establish

$$(1-\delta)\|x\|_2^2 \leq \|\Phi x\|_2^2 \leq (1+\delta)\|x\|_2^2$$

"the easy way"

• Re-write the RIP as a the supremum of a random process

$$\sup_{x} |G(x)| = \sup_{x} |x^* \Phi^* \Phi x - x^* x| \le \delta$$

where the sup is taken over all  $2S\mbox{-sparse}$  signals

• Estimate this sup using tools from probability theory (e.g. the Dudley inequality) — approach pioneered by Rudelson and Vershynin

# Random convolution

 Many active imaging systems measure a pulse convolved with a reflectivity profile (Green's function)



- Applications include:
  - radar imaging
  - sonar imaging
  - seismic exploration
  - channel estimation for communications
  - super-resolved imaging
- Using a random pulse = compressive sampling (Tropp et al. '06, R '08, Herman et al. '08, Haupt et al. '09, Rauhut '09)

# Random convolution for CS, theory

- $\bullet$  Signal model: sparsity in any orthobasis  $\Psi$
- Acquisition model: generate a "pulse" whose FFT is a sequence of random phases (unit magnitude),

convolve with signal,

sample result at m random locations  $\Omega$ 

$$\Phi = R_{\Omega} \mathcal{F}^* \Sigma \mathcal{F}, \quad \Sigma = \operatorname{diag}(\{\sigma_{\omega}\})$$

• The RIP holds for (R '08)

$$M \gtrsim S \log^5 N$$

Note that this result is *universal* 

• Both the random sampling and the flat Fourier transform are needed for universality

# Randomizing the phase



sample here

# Why is random convolution + subsampling universal?

• One entry of  $M = \mathcal{F}\Sigma\hat{\Psi}$ :

$$M_{t,s} = \sum_{\omega} e^{j2\pi\omega t} \sigma_{\omega} \hat{\psi}_s(\omega)$$
$$= \sum_{\omega} \sigma'_{\omega} \hat{\psi}_s(\omega)$$

• Size of each entry will be concentrated around  $\|\hat{\psi}_s(\omega)\|_2 = 1$ does not depend on the "shape" of  $\hat{\psi}_s(\omega)$ 

# Compare to Fast Johnson-Lindenstrauss Transform

- Ailon and Chazelle, 2006
- Problem:

k points  $x_1, \ldots, x_k$  in  $\mathbb{R}^N$ , project onto  $\mathbb{R}^M$  using  $\Phi$  ( $M \times N$  matrix) Want  $\|\Phi(x_i - x_j)\|_2 \approx \|x_i - x_j\|_2$  for  $M \sim \log k$ , and  $\Phi$  to be "fast"

- JL problem is closely related to CS (Baraniuk et al. '07)
- Their solution: take  $\Phi = PHD$   $D = \text{diag}(\{\epsilon_i\})$  (makes input signs random) H = Hadamard transform (Fourier on  $\mathbb{Z}_2$ )  $P = M \times N$  subsampling matrix, each row has m random entries at random locations
- This  $\Phi$  would be tremendous, except it is not clear how to implement it by taking O(M) physical measurements (*P* has  $M^2$  entries in it)

# Seismic forward modeling

- Run a single simulation with all of the sources activated simultaneously with random waveforms
- The channel responses interfere with one another, but the randomness "codes" them in such a way that they can be separated later



Related work: Herrmann et. al '09

### Restricted isometries for multichannel systems



• With each of the pulses as iid Gaussian sequences,  $\Phi$  obeys

$$\begin{split} (1-\delta)\|h\|^2 &\leq \|\Phi h\|_2^2 &\leq (1+\delta)\|h\|_2^2 \quad \forall s\text{-sparse } h\in \mathbb{R}^{nc} \\ \text{when} & (\text{R and Neelamani '09}) \end{split}$$

$$m \gtrsim s \cdot \log^5(nc) + n$$

• **Consequence:** we can separate the channels using short random pulses (using  $\ell_1$  min or other sparse recovery algorithms)

# Summary

• Main message of CS:

We can recover an S-sparse signal from  $\sim S\log N$  measurements

- Random matrices (iid entries)
  - easy to analyze, optimal bounds
  - univeral
  - hard to implement and compute with
- Structured random matrices (random sampling, random convolution)
  - structured, and so computationally efficient
  - physical
  - much harder to analyze, bound with exta log-factors