

# Sublinear Time and Space Algorithms 2026A – Problem Set 5

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**General instructions:** Please keep your answers short and easy to read. You can use results, calculations or notation seen in class without repeating them, unless asked explicitly to redo them.

1. Design a deterministic algorithm that given a set of vertices  $S \subseteq V$  finds a neighbor of it using  $O(n^{1/r})$  queries in  $O(r)$  rounds. Then, use this to create a deterministic algorithm that determines if a graph is connected using  $O(n^{1+1/r})$  queries in  $O(r \log n)$  rounds.
2. Design a sublinear-time algorithm that, given  $0 < \varepsilon < 1$  and access to an  $n$ -vertex graph  $G$ , approximates the number of connected components, within additive error  $\varepsilon n$ . The access to  $G$  allows to query, for each vertex, its degree and its neighbors list (e.g., ask for the  $i$ -th neighbor). What is the running time of your algorithm?

Hint: Start with the special case that all connected components have size at most  $1/\varepsilon$ . Consider sampling a uniformly random vertex  $v$  and computing its connected component  $C_v$ .